

Geometric Simplicity Theory

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T will be a complete theory in \mathcal{L} having only infinite models.
 We work in a big saturated model \mathcal{M} of size $\bar{\kappa}$ of T .
 M, N, \dots will be (small) elementary submodels of \mathcal{M} ;
 A, B, \dots will be subsets of \mathcal{M} ;
 a, b, c, \dots denote tuples (possibly infinite) from \mathcal{M} , and $c \in A$
 means c is a tuple from A ;
 AB , or Ac with $c = (c_1 \dots c_n)$ often abbreviates $A \cup B$, or
 $A \cup \{c_1, \dots, c_n\}$ respectively;
 $a \equiv_A b$ means $\text{tp}(a/A) = \text{tp}(b/A)$.
 Types $p(x), q(x), \dots$ are partial types with parameter, so are
 formulas $\varphi(x), \psi(x), \dots$. For $\varphi(x, a)$, a is the parameter on
 $\varphi(x, y) \in \mathcal{L}$.

Definition

- 1
 - A formula $\varphi(x, a_0)$ *k-divides over A* if there is an A -indiscernible sequence $\langle a_i \mid i < \omega \rangle$ such that $\{\varphi(x, a_i) \mid i < \omega\}$ is k -inconsistent.
 - A formula $\varphi(x, a)$ *divides over A* if it k -divides for some positive k .
 - A type $p(x)$ *divides over A* if $p \vdash \varphi(x, a)$ and $\varphi(x, a)$ divides over A .
- 2 A type $p(x)$ *forks over A* if $p \vdash \varphi_0(x, a_0) \vee \dots \vee \varphi_k(x, a_k)$ and $\varphi_i(x, a_i)$ divides over A for each $i \leq k$.

Write $A \downarrow_B C$ if for any $a \in A$, $\text{tp}(a/BC)$ does not fork over B .

- (Extension) Let $A \subseteq B \subseteq C$. If $p \in S(B)$ does not fork over A , then there is an extension $q \in S(C)$ of p , which does not fork over A ; equivalently, if $c \downarrow_A B$, then there is $c' \equiv_B c$ such that $c' \downarrow_A C$.

Fundamental Theorem of Forking

The following are equivalent.

- (Symmetry) $A \downarrow_B C$ iff $C \downarrow_B A$.
- (Transitivity) For $A \subseteq B \subseteq C$ and d , $d \downarrow_A C$ iff $d \downarrow_A B$ and $d \downarrow_B C$.
- (Local Character) For any set A and finite d , there is $A_d (\subseteq A)$ of size $\leq |\mathcal{L}|$ such that $d \downarrow_{A_d} A$.

Definition

- T is said to be *simple* if one of the equivalent properties of forking holds.
- T is *unstable* if there is $\varphi(x, y)$ and $a_i, b_i \in \mathcal{M}$ ($i < \omega$) such that $\mathcal{M} \models \varphi(a_i, b_j)$ iff $i < j$.
- T is *stable* if it is not unstable.

$\mathcal{C} = \{C_0, C_1, \dots\}$ is A -independent iff $C_{i+1} \downarrow_A C_0 \dots C_i$ for all i
iff $C_i \downarrow_A \mathcal{C} - \{C_i\}$ for all i .

Fact

- If T is stable, then it is simple.
- If T is simple, then forking is dividing. Hence $A \downarrow_B C$ iff for any $c \in A$, $\text{tp}(c/BC)$ does not divide over A .

uncountable categorical (ACF) \subseteq superstable (DCF) \subseteq stable (SCF). Ordered field is unstable.

superstable \subseteq supersimple (PsF) \subseteq simple (PAC)
 \subseteq stable

Fact

(Shelah)

- If T is stable, then \downarrow additionally satisfies the following axiom.

(Uniqueness over a model) For $M \subseteq A$, if $c \equiv_M c'$ and $c \downarrow_M A$, $c' \downarrow_M A$, then $c \equiv_A c'$.

- The four basic axioms together with above axiom characterize stability and forking. Namely, if there is an invariant relation \downarrow^* between tuple and sets satisfying the 5 axioms, then T is stable and $\downarrow^* = \downarrow$.

Question

For simplicity, what axiom can substitute for (Uniqueness over a model) ?

Theorem

(K, Pillay)

- If T is simple then \perp additionally satisfies the following axiom.

(The Independence Theorem (or Type amalgamation) over a model) For $M \subseteq A_0, A_1$, if $c_0 \equiv_M c_1$, $A_0 \perp_M A_1$, and $c_i \perp_M A_i$ ($i = 0, 1$), then there is $c \equiv_{A_i} c_i$ such that $c \perp_M A_0 A_1$.

- The four basic axioms together with above axiom characterize simplicity and forking.

Definition

- We say a, b have the same strong type over A (write $a \equiv_A^s b$, or $\text{stp}(a/A) = \text{stp}(b/A)$) if $E(a, b)$ holds, for any finite definable equivalence relation $E(x, y)$ over A .
- We say a, b have the same Lascar (strong) type over A (write $a \equiv_A^L b$, or $\text{Ltp}(a/A) = \text{Ltp}(b/A)$, or $\text{Lstp}(a/A) = \text{Lstp}(b/A)$) if $\mathbf{E}(a, b)$ holds, for any bounded type-definable equivalence relation $\mathbf{E}(x, y)$ over A .

Theorem

- Let T be stable.

(Uniqueness for strong types) For $B \subseteq C$, if $a \equiv_B^s a'$ and $a \perp_B C$, $a' \perp_B C$, then $a \equiv_C a'$.

- Let T be simple.

(Type amalgamation for Lascar types) For $B \subseteq C_0, C_1$, if $a_0 \equiv_B^L a_1$, $C_0 \perp_B C_1$, and $a_i \perp_B C_i$ ($i = 0, 1$), then there is $a \equiv_{C_i} a_i$ such that $a \perp_B C_0 C_1$.

Open Problem

In simple T , does type amalgamation hold for strong types ?

Definition

- ① For any type $p(x)$, $D(p(x), \varphi(x, y), k)$ rank is defined by induction as follows:

 - $D(p, \varphi, k) \geq 0$ for any consistent type p .
 - $D(p, \varphi, k) \geq n + 1$ if for some a , $\varphi(x, a)$ k -divides over $\text{dom}(p)$, and $D(p \cup \{\varphi(x; a)\}, \varphi, k) \geq n$.
- ② $D(p(x), \varphi(x, y))$ rank is defined as: $D(p, \varphi) \geq \alpha + 1$ if for some a , $\varphi(x, a)$ divides over $\text{dom}(p)$, and $D(p \cup \{\varphi(x; a)\}, \varphi) \geq \alpha$.
- ③ $D(\psi(x))$ rank is defined as: $D(\psi(x)) \geq \alpha + 1$ if there is a formula $\varphi(x)$ dividing over $\text{dom}(\psi(x))$, and $D(\psi(x) \wedge \varphi(x)) \geq \alpha$;
 $D(p(x)) := \min\{D(\psi_1(x) \wedge \dots \wedge \psi_n(x)) \mid \psi_i(x) \in p(x)\}$.

Fact

The following are equivalent.

- *T has the tree property, i.e. there exist a formula $\varphi(x, y)$, an integer $k \geq 2$ and tuples c_α with $\alpha \in \omega^{<\omega}$ such that*
 - (i) for any $\alpha \in \omega^{<\omega}$, $\{\varphi(x, c_{\alpha \frown n}) \mid n \in \omega\}$ is k -inconsistent,*
 - (ii) for all $\beta \in \omega^\omega$, $\{\varphi(x, c_{\beta \upharpoonright n}) \mid n \in \omega\}$ is consistent.*
- *$D(x = x, \varphi, k) = \infty$ for some $\varphi, k \geq 2$.*

Fact

- *The following are equivalent.*
 - *T is simple.*
 - *$D(x = x, \varphi, k) < \infty$ for any φ, k .*
 - *$D(x = x, \varphi, k) < \omega$ for any φ, k .*
- *$D(x = x, \varphi, k) \leq D(x = x, \varphi) \leq D(x = x)$.*

Definition

Simple T is said to be *low* if $D(x = x, \varphi) < \omega$ for any φ .

Fact

- T is low iff for any $\varphi(x, y)$, there is k such that whenever $\varphi(x, a)$ divides, it k -divides.
- (S. Buechler) If T is low, then Lascar types are strong types (i.e. $a \equiv_A^s b$ iff $a \equiv_A^L b$). Equivalently, type amalgamation holds for strong types.

Definition

- T is said to be *supersimple* if for any finite c and A , there is finite $A_c \subseteq A$ such that $c \perp_{A_c} A$.
- T is *superstable* if T is stable and supersimple.

Definition

For $p \in S(A)$, $SU(p) \geq \alpha + 1$ if if p has a forking complete extension q s.t. $SU(q) \geq \alpha$.

Fact

The following are equivalent.

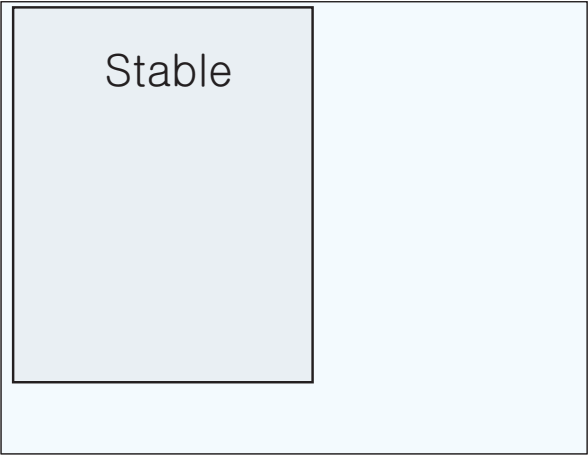
- T is supersimple.
- $D(x = x) < \infty$.
- $SU(p) < \infty$ for any complete p .

Fact

- T simple. $c \downarrow_A B$ iff $D(\text{tp}(c/A), \varphi, k) = D(\text{tp}(c/AB), \varphi, k)$ for any φ, k .
- T supersimple. $c \downarrow_A B$ iff $SU(c/AB) = SU(c/A)$. (But D -rank reflects forking only for superstable T .)

Theorem

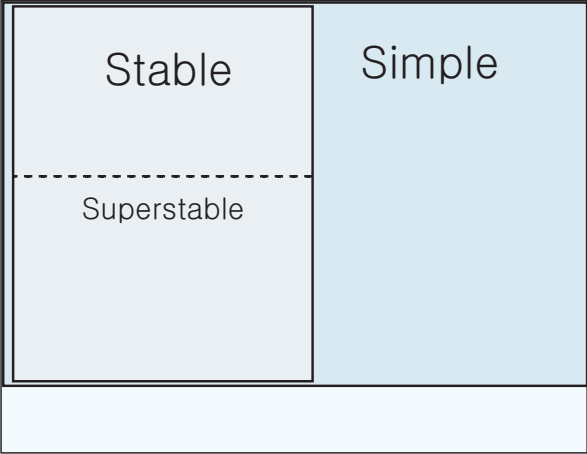
- (Lachlan) If T is countable superstable, then either $I(T, \omega) = 1$ or $\geq \omega$.
- (K) If T is countable supersimple, then either $I(T, \omega) = 1$ or $\geq \omega$.

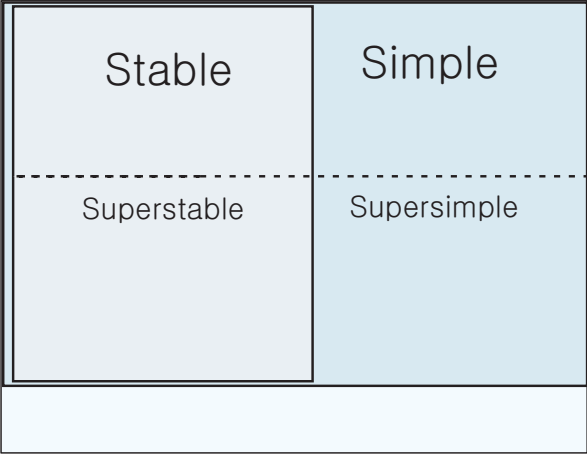


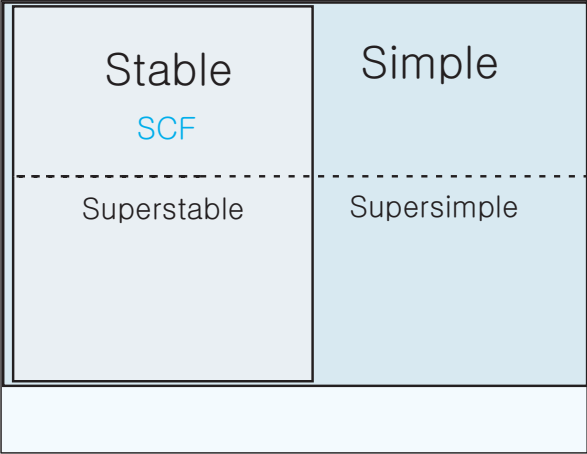
Stable

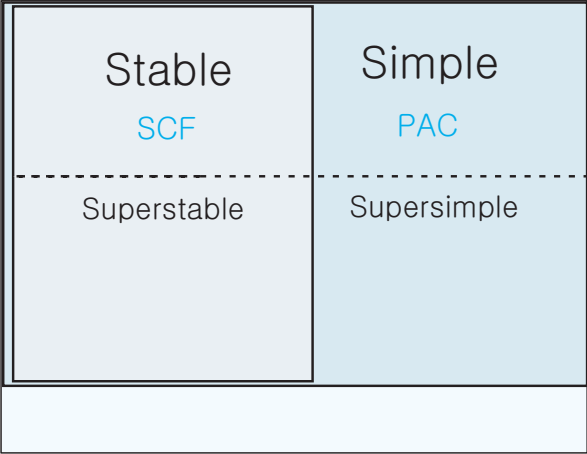
Stable

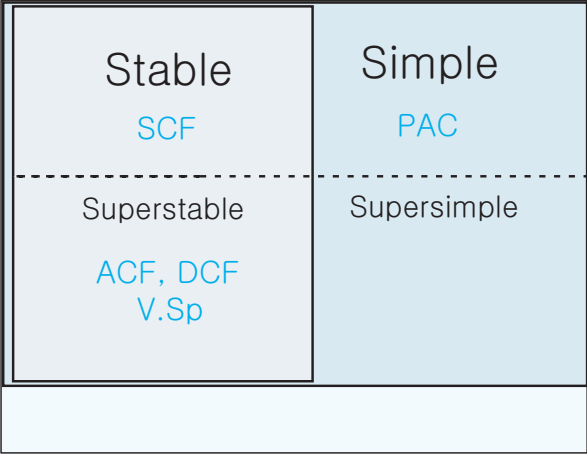
Simple



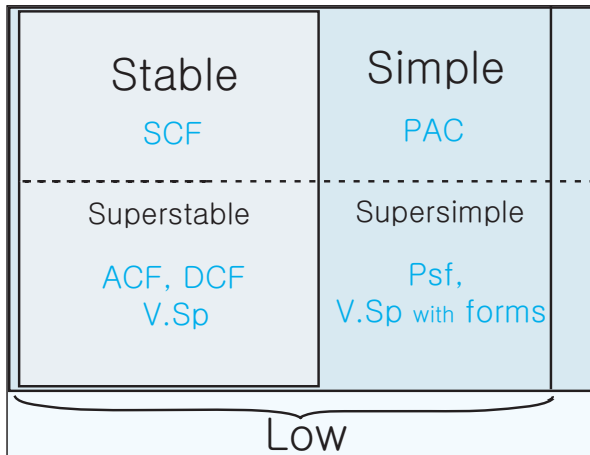








Stable SCF	Simple PAC
Superstable ACF, DCF V.Sp	Supersimple Psf, V.Sp with forms



Definition

- T has *elimination of imaginaries* if for any definable equivalence class a/E , there is finite $b \in \mathcal{M}$ such that they are interdefinable (i.e. for any automorphism fixes the class a/E setwise iff it fixes the tuple b).
- T has *elimination of hyperimaginaries* if for any type-definable equivalence class \bar{a}/\mathbf{E} (\bar{a} possibly infinite), there are definable equivalence classes $\{a_i/E_i\}_i$ such that they are interdefinable.

From now on, we assume T has elimination of imaginaries, and elimination of hyperimaginaries, (so assume $Ltp = stp$).

Unless mentioned otherwise T will be simple.

Theorem

(Buechler, Pillay, Wagner) If T is supersimple, then T has elimination of hyperimaginaries.

Fact

If type amalgamation holds for $tp(a/A)$, then indeed $tp(a/A) = Ltp(a/A) = tp(a/\text{acl}(A))$. Hence type amalgamation holds for any acl -set (e.g. a model). Such a type we call an amalgamation base.

Theorem

(Hart, K, Pillay) Any amalgamation base has a canonical base. Namely, for any amalgamation base $p = \text{tp}(c/A)$ with $A = \text{acl}(A)$, there is smallest $A_0 \subseteq A$ such that $c \perp_{A_0} A$ and $\text{tp}(c/A_0)$ is an amalgamation base. We write $A_0 := \text{Cb}(c/A)$.

Observation

Let $B_i = \text{acl}(B_i) \subseteq B$ ($i = 0, 1$). If $A \perp_{B_0} B$ and $A \perp_{B_1} B$, then $A \perp_{B_0 \cap B_1} B$: Note that $C := \text{Cb}(A/B) \subseteq B_i$, \therefore by transitivity, $A \perp_C B_0 \cap B_1$ and $A \perp_{B_0 \cap B_1} B$ hold.

Definition

Let $p(x) \in S(A)$ be an amalgamation base. We call $\varphi(x, y)$, *p-stable* if the φ -type $p \upharpoonright \varphi$ has a unique nonforking φ -type extension over any $B \supseteq A$.

Fact

- $\varphi(x, y)$ is stable iff $\varphi(x, y)$ is *p-stable* for any amalgamation base $p(x)$.
- If $\varphi(x, y)$ is *p-stable*, then there is $\psi_\varphi(y)$ over $\text{Cb}(p)$, called the φ -definition of p , such that $\models \psi_\varphi(b)$ holds iff $\varphi(x, b)$ is in some nonforking extension of p .

Theorem

(K) Assume T is supersimple, and $p(x)$ is an amalgamation base.
Then

$\text{Cb}(p) = \{ \text{names of } \psi_\varphi(y) \mid \varphi(x, y) \text{ is } p\text{-stable} \}$.

Definition

By a *Morley sequence* of $p \in S(A)$, we mean an A -indiscernible and independent sequence $I = \langle a_i \mid i < \omega \rangle$ in p .

$\text{Cb}(p) \subseteq I$ where I is a Morley sequence of p .

Consider a vector space $V = (V, +, r)_{r \in F}$. For acl-subsets A, B of V , we have $A \downarrow_{A \cap B} B$. Equivalently, when finite dimensional, $\dim(A \cup B) = \dim(A) + \dim(B) - \dim(A \cap B)$.

Definition

- T is called *modular* if \downarrow satisfies above relations between acl-sets.
- T is said to be *1-based* if for any indiscernible sequence $I = \langle a_i \mid i < \omega \rangle$, $I - a_0$ is Morley over a_0 .

Observation

The following are equivalent.

- (i) T is 1-based.
- (ii) For any $p = \text{tp}(c_0/A)$, $\text{Cb}(c_0/A) \subseteq \text{acl}(c_0)$.
- (iii) T is modular.

Proof: (i) \Rightarrow (ii) Let $e := \text{Cb}(c_0/A)$. Now there is a Morley sequence $J = \langle c_i \mid i \leq \omega \rangle$ of p . Then $c_\omega \perp_A I$ where $I = J - \{c_\omega\}$. Since $c_\omega \perp_e AI$ and $e \subseteq I$, we have $c_\omega \perp_I A$. Hence $e = \text{Cb}(c_\omega/A) = \text{Cb}(c_\omega/I)$. By 1-basedness, $c_\omega \perp_{c_0} I$. Hence $e \in \text{acl}(c_0)$.

(ii) \Rightarrow (iii) Given A, B , we have $A \perp_{\text{Cb}(A/B)} B$ and $\text{Cb}(A/B) \subseteq \text{acl}(A)$. Hence by transitivity, $A \perp_{\text{acl}(A) \cap \text{acl}(B)} B$.

(iii) \Rightarrow (i) Indiscernible $I = \langle c_i \mid i < \omega \rangle$ is given. (iii) implies $c_0 \downarrow_A c_1 c_2 \dots$ where $A = \text{acl}(c_0) \cap \text{acl}(c_1 c_2 \dots)$. Now, $tp(c_1 c_2 \dots / \text{acl}(c_0)) = tp(c_2 c_3 \dots / \text{acl}(c_0))$. Hence $A = \text{acl}(c_0) \cap \text{acl}(c_2 c_3 \dots)$. Again, $tp(c_0 / \text{acl}(c_2 c_3 \dots)) = tp(c_1 / \text{acl}(c_2 c_3 \dots))$, and thus $A = \text{acl}(c_1) \cap \text{acl}(c_2 c_3 \dots)$. Then, due to (iii) again, $c_1 \downarrow_A c_2 c_3 \dots$. From transitivity, we have $c_1 \downarrow_{c_0 A} c_2 c_3 \dots$. Then $c_1 \downarrow_{c_0} c_2 c_3 \dots$ since $A \subseteq \text{acl}(c_0)$. Thus (i) follows.